Modeling UNO Gameplay Using Graph Theory to Analyze Optimal Strategies

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Abstract— This paper presents a graph theory-based approach to analyze and optimize strategies in the card game UNO. We model the gameplay using weighted graphs and implement both Dijkstra's and Minimax algorithms for strategic decision-making. Our implementation demonstrates that Dijkstra's algorithm excels at optimizing immediate moves, while Minimax provides robust defensive strategies in multi-player scenarios. The results show that graph-based modeling effectively identifies optimal card sequences and strategic decisions, contributing to the broader understanding of applying graph theory in game strategy optimization.

Keywords—graph theory, UNO card game, Dijkstra's algorithm, Minimax algorithm, game strategy optimization.

I. INTRODUCTION

Graph theory, a branch of discrete mathematics, is widely used to model relationships and optimize decision-making processes in complex systems. It provides a structured approach to understanding interactions within networks, making it a powerful tool for solving problems in various fields such as computer science, logistics, and social networks. The application of graph theory to games, particularly those involving strategic decision-making, has gained significant attention due to its ability to uncover optimal strategies and predict outcomes.

UNO, a popular card game with simple yet dynamic rules, presents a unique opportunity for applying graph theory in a recreational setting. In UNO, each card is characterized by its color, number, and specific actions (e.g., skip, reverse, draw two), and players are required to match cards based on color or number to make moves. This leads to a system where the relationships between cards and players can be effectively modeled using graph theory.



Fig 1. Uno Games

Source: Playstation Store App

In this paper, we model the game of UNO using graphs, where each card is represented as a vertex and edges connect cards that can be played consecutively according to the game's rules. The graph will help identify optimal strategies by analyzing the connections between cards and their effects on gameplay. Specifically, we explore how the structure of the graph—such as the adjacency of playable cards, the potential sequences of moves, and the strategic positioning of action cards—can inform decision-making.

Through this approach, we aim to provide a deeper understanding of the strategic elements in UNO, answering questions such as:

- What is the most efficient sequence of moves to reduce hand size quickly?
- How can players identify which cards should be prioritized to maintain a playable hand?
- How does the presence of action cards (e.g., skip, reverse) influence the optimal strategy?

By modeling the gameplay through a graph, we not only highlight the potential of graph theory in game theory but also demonstrate how it can be applied to enhance decision-making in simple rule-based environments. The results of this study offer valuable insights into strategy optimization, with broader implications for games involving similar structured interactions.

II. THEORETICAL BASIS

1. Graph Theory

Graph theory, a pivotal area of discrete mathematics, examines the relationships between pairs of objects through graphs. A graph comprises **vertices** (nodes) and **edges** (connections between nodes). In this study, each **vertex** represents a card in the UNO game, and each **edge** signifies a playable connection between two cards, adhering to the game's rules [1].

Graphs can be categorized into various types, including:

- **Bipartite Graphs**: Graphs where vertices can be divided into two disjoint subsets, with edges only connecting vertices from different subsets. In UNO, this could represent cards classified by color or number [1].



Fig 2.1. Bipartite Graph Source: Graph-Part-1 informatika.stei.itb.ac.id

- Weighted Graphs: Graphs where edges have weights assigned, representing some form of cost or distance. In UNO, weights could denote the strategic value of a card in terms of its impact on the game's progress [1].



Fig 2.2. Weighted Graph and Unweighted Graph Source: Graph-Part 1 informatika.stei.itb.ac.id

- **Trees and Cycles:** A tree is a connected graph with no cycles, while a cycle is a path where the starting and ending vertices are the same. Cycles may appear in the sequence of moves during the game, where cards loop back to earlier positions [1].



Fig 2.3 Tree, Cycles, and Forest Source: Graph-Part 1 informatika.stei.itb.ac.id

2. Graph Representation of UNO

In this study, we model the UNO game using a graph where each card is represented by a vertex, and edges connect cards that are playable based on the game's rules. The game follows a set of simple yet dynamic rules:

- Players can play a card if it matches the color or

number of the card in the centre [3].

Special action cards (such as **skip**, **reverse**, **draw two**, **wild**) have additional effects that modify the sequence of moves or the flow of the game [3].

Thus, in our graph representation:

- Vertices represent the cards in a player's hand or the deck.
- **Edges** represent possible moves between cards that satisfy the game's playability rules (matching colors or numbers).
- Action cards introduce additional edges that affect the state of the game (such as skipping the next player or reversing the direction of play).

3. Graph Algorithms for Strategy Analysis

Several key graph algorithms can be applied to analyze optimal strategies in UNO:

3.1. **Depth-First Search (DFS) and Breadth-First Search (BFS)**: DFS explores the graph by diving deep into one possible sequence of moves before backtracking. This can be used to simulate a sequence of card plays to find the longest playable sequence or to identify dead ends in gameplay. By exploring all possible paths, DFS can help analyze which cards should be played first to maximize future opportunities [4].): BFS explores the graph level by level, ensuring that all reachable cards are explored in a systematic order. In the context of UNO, BFS could help identify the shortest path (i.e., the fewest number of moves) to reach a winning state by playing all cards [4].



Fig 2.4. DFS and BFS Representation Source: GeeksforGeeks

- 3.2. **Maximum Matching**: This algorithm helps identify pairs of cards that can be played consecutively based on the game's rules. By maximizing the number of matching cards in the player's hand, we can strategize the optimal sequence for card play [4].
- 3.3. Graph Coloring: Graph coloring can be used to represent situations where cards of a certain color must be prioritized or separated. The coloring algorithm can help in situations where multiple color cards exist, and finding an optimal color grouping for a player's hand can be beneficial [4].



Fig 2.5. Graph Coloring Source: Graph-Part 3-informatika.stei.itb.ac.id

- 3.4. Shortest Path Algorithms (Dijkstra's or Bellman-Ford): These algorithms can be used to calculate the quickest path to reduce a player's hand size. In the context of UNO, the goal is to minimize the number of turns or moves required to empty the hand, taking into account action cards that could alter the course of play [4].
- 3.5. **Minimax Algorithm**: The Minimax Algorithm is a decision-making strategy typically used in two-player games where one player's gain is another player's loss. In the context of UNO, the algorithm can be adapted to model strategic decision-making in dynamic, multiplayer scenarios by treating the game's state as a tree structure, where each node represents a possible game state, and edges represent potential moves.

4. Strategy Optimization in Game Theory

In game theory, the concept of **optimal strategies** is crucial. An optimal strategy is one that maximizes a player's chances of winning, taking into account all possible moves and counter-moves. The optimal strategy in UNO is influenced by both the state of the game (i.e., the cards in the hand and on the table) and the rules governing card play [5].

In the context of UNO, graph theory helps optimize strategy by:

- **Modeling the possible sequences of plays** to minimize the number of moves required to win [5].
- **Analyzing the effects of action cards** (such as Wild and Skip) to determine the best time to use them for disrupting the opponent's progress [5].
- **Identifying the best starting card** by analyzing the graph to determine which card is most likely to open up the greatest number of future plays [5].

By combining these graph-based algorithms with gametheory concepts, it is possible to develop a comprehensive strategy that can be applied in real-world UNO gameplay [5].

III. ANALYZING UNO STRATEGY USING GRAPH-BASED APPROACH

Based on the theorical basis that had been discussed, as the topic of this paper is bound to a strategic games. The graph is

used to represents not only the standard UNO card that used in the game, but also the game-flow is represented by using graph. The graph that is being used is basically to generate the best algorithm procedure to find the best move it could find. For the example in this paper is going to used the Dijkstra's Algorithm (an algorithm to find the shortest path by using weighted value by weighted graph represented).

1. Cards Value Representation Using Graph

In the context of UNO, each card possesses unique attributes such as color, number, or action type (e.g., Skip, Reverse, Wild). These attributes can be effectively mapped into a graph structure, where each card is represented as a **vertex** and their playable relationships are defined as **edges**. The edges in the graph signify the possible moves allowed based on UNO rules, such as matching by color, number, or action.

1.1. Vertex Representation

Each vertex in the graph corresponds to a specific card in the UNO deck. For instance: A red "5" card is represented as a vertex labelled **R5**. A blue "Skip" card is represented as a vertex labelled **B-Skip**.

1.2. Edges Representation

Edges between vertices define the transitions between playable cards:

- O If two cards share the same color, an edge is formed. For example: R5 ↔ R7 (Red cards "5" and "7").
- If two cards share the same number, an edge is formed. For example: R5 ↔ G5 (Cards "5" in red and green colors).
- Action cards (e.g., Wild) connect to all possible playable cards. For example: Wild ↔ R5, Wild ↔ B-Skip.

1.3. Weighted Graph Representation

Edges between vertices define the transitions between playable cards:

- O If two cards share the same color, an edge is formed. For example: R5 ↔ R7 (Red cards "5" and "7").
- If two cards share the same number, an edge is formed. For example: R5 ↔ G5 (Cards "5" in red and green colors).
- Action cards (e.g., Wild) connect to all possible playable cards. For example: Wild ↔ R5, Wild ↔ B-Skip.



Fig 3.1. Graph for cards value representation Source: Author's Document

2. Game Flow Representation Using Graph

The game flow in UNO can also be represented as a dynamic graph. Here, the focus is on the sequence of moves during a player's turn and how the graph evolves with each play.

2.1. Turn-Based Transitions

During each turn, the graph is updated based on the active card on the discard pile:

- The active card becomes the "source" vertex.
- All valid moves from the player's hand form the outgoing edges from the source vertex.
- 2.2. Example of Turn-Based Graph Evolution Suppose the active card is **R5**, and the player's hand contains {R7, G5, Wild}:
 - Graph at turn start: Source vertex: **R5**. Outgoing edges: $\{R5 \rightarrow R7, R5 \rightarrow G5, R5 \rightarrow Wild\}$.
 - o After playing **R7**, the graph updates: Source vertex: **R7**. Outgoing edges: {R7 → Wild} (assuming only Wild remains playable).
- 2.3. Dynamic Weight Adjustment

To optimize the decision-making process, edge weights can dynamically adjust based on:

- Remaining cards: Assign higher weights to moves that leave the opponent with more cards.
- Action cards: Prioritize moves that disrupt opponents (e.g., Skip, Reverse).



Fig 3.2. Graph for game-flow Source: Author's document

- 3. Strategic Pathfinding Using Dijkstra's Algorithm The graph structure allows for the application of Dijkstra's algorithm to find the optimal sequence of plays. The goal is to identify the shortest path (lowest weight) to achieve a winning condition, such as emptying the hand.
 - 3.1. Algorithm Implementation
 - **Input Graph**: Provide the weighted graph representing the current state of the game.
 - o Start Node: Set the active card as the starting

vertex.

- End Node: Define the goal state, such as a vertex representing an empty hand or a specific winning move.
- **Pathfinding**: Apply Dijkstra's algorithm to calculate the shortest path and its total weight.
- 3.2. Example Application
 - Initial state: Player's hand = {R7, G5, Wild}, Active card = R5.
 - Graph weights:
 - $R5 \rightarrow R7$: Weight 1 (matching color).
 - $R5 \rightarrow G5$: Weight 2 (matching number).
 - R5 \rightarrow Wild: Weight 0.5 (universal playability).
 - Dijkstra's output:
 - Optimal path: $R5 \rightarrow Wild$.
 - Total weight: 0.5.



Fig 3.3. Dijkstra's algorithm pathfinding graph visualization Source: Author's document

Move	Weight (w)	Effectiveness $(\frac{1}{w})$	Percentage $(\frac{E}{\Sigma E} \times 100\%)$
R7 (Same	1	1	28,57%
Color)			
G5 (Same	2	0,5	14,29%
Number)			
Wild	0,5	2	57,14%
(Universal)			

4. Dynamic Strategy Optimization Using Minimax Algorithm

4.1. Representing the Game Using a Game Tree

To apply the Minimax Algorithm, the game is first represented as a tree structure. Each node in the tree corresponds to a unique game state, defined by:

The player's hand (cards available for play).

- The current top card on the discard pile.
- The cards potentially held by the opponents (estimated based on previous moves).
- The state of the draw pile.

Edges between nodes represent possible actions, such as playing a card, drawing a card, or skipping a turn. The game tree is expanded up to a specific depth, representing a sequence of moves. Terminal nodes (leaf nodes) correspond to the end of a series of moves, where the evaluation function assigns a score to the resulting game state.

4.2. Evaluation Function

The evaluation function is central to the Minimax Algorithm, as it quantifies the desirability of each game state. For UNO, the evaluation function could consider:

- **Reduction in Player's Cards**: Favor moves that reduce the number of cards in the player's hand.
- **Impact on Opponents**: Penalize moves that allow opponents to play or reduce their cards easily.
- Special Card Effects: Reward moves that utilize special cards (e.g., Skip, Reverse, Wild) to disrupt opponents or gain strategic advantage.
- **Game Progression**: Account for moves that bring the player closer to winning the game while limiting the opponents' ability to do so.

4.3. Applying the Minimax Algorithm

The Minimax Algorithm is applied by first constructing a game tree from the current state of the UNO game, where each branch represents a sequence of possible moves. Starting at the root node, which corresponds to the current game state, the algorithm explores all potential actions up to a specific depth in the tree. Each terminal node (leaf node) is evaluated using an evaluation function that assigns a score to the resulting game state. The algorithm then backpropagates these scores to determine the optimal path. Maximizing nodes, which represent the current player's turn, select the move that yields the highest score, aiming to improve the player's advantage. Conversely, minimizing nodes, representing the opponents' turns, choose the move with the lowest score, attempting to minimize the maximizing player's benefit. This process ensures that the selected move considers not only the player's perspective but also the opponents' potential responses. Additionally, after each turn, the game tree is dynamically recalculated, allowing the strategy to adapt to the evolving game state and maintain its effectiveness throughout the game.

4.4. Scenario Example

Consider a scenario where the player's hand includes the following cards: R5, R7, G5, B-Skip, Wild. The current top card on the pile is R5. The possible moves are:

- Play R7 (same color as the top card).
- Play G5 (same number as the top card).
- Play Wild (valid regardless of the top card).
- Play B-Skip (invalid unless a blue card or Skip is on top, so not considered).

o Draw a card from the draw pile.



Fig 4.1. Visualization for game scenario examples Source: Author's document

The Minimax Algorithm evaluates these moves, considering the effects on the game state and opponents. For instance, playing Wild might force the opponent to draw cards if they lack the chosen color, while playing R7 maintains the current color and avoids helping the opponent.

4.5. Algorithm Implementation

The implementation of the Minimax Algorithm in the context of UNO requires constructing a game tree, where each node represents a game state. The edges denote transitions resulting from possible moves, and terminal nodes are evaluated to determine the desirability of a game state. The algorithm backpropagates these evaluations to identify the optimal move.

import math		
import networkx as nx		
import matplotlib.pyplot as plt		
# Example cards and evaluation values		
card_values = {		
'R5': 5,		
'R7': 7,		
'Wild': 10.		
'B-Skip': 3,		
'end': 0		
}		
# Example game tree		
game_tree = {		
'start': ['R5', 'Wild'],		
'R5': ['R7', 'G5'], 'Wild': ['P.Svip' 'P7']		
'R7': ['end'].		
'GS': ['end'],		
'B-Skip': ['end']		
}		
<pre># Minimax Algorithm with Graph Coloring def minimax(state, depth, maximizing_player, graph=None, parent=None):</pre>		
unu Talaasha bha Wisiasu Alaashba sidh aashb alaaisa		
implements the minimax Algorithm with graph coloring.		
if graph is not None:		
# Assign node color based on player type or terminal state		
if state not in game_tree: # Terminal state		
else:		
<pre>color = "green" if maximizing_player else "red"</pre>		
<pre>graph.add_node(state, value=card_values.get(state, -math.inf), color=color)</pre>		
if parent is not None:		
grapn.add_edge(parent, state)		
# Base case: terminal state or depth limit		
if state not in game_tree or depth == 0:		
return card_values.get(state, -math.inf)		
if maximizing_player:		
<pre>max_eval = -math.inf</pre>		
for child in game_tree[state]:		
max eval = max(max eval, eval)		
return max_eval		
min_eval = math.inf		
eval = minimax(child, depth - 1, True, graph, state)		
min_eval = min(min_eval, eval)		
return min_eval		

Fig 4.2. Source code for minimax algorithm (based on the example given) Source: Author's Document

The **Minimax Algorithm** implemented in the code is designed to evaluate the best move in a game scenario, represented as a tree structure. The algorithm recursively explores all possible moves from a given state, alternating between a maximizing player (who seeks to maximize the score) and a minimizing player (who seeks to minimize the score).

In the provided example, the game starts with the "start" state. From here, the player can choose between playing the card **R5** (Red 5) or **Wild**. If the player chooses **R5**, the subsequent moves could be **R7** (Red 7) or **G5** (Green 5). If the player chooses **Wild**, the next moves could be **B-Skip** (Blue Skip) or **R7**. Each of these moves leads to an **end** state, which signifies the termination of the game.

The algorithm begins at the root node ("start") and evaluates each branch of the tree to a specified depth (3 in this case). For instance, if the maximizing player selects **R5**, the algorithm proceeds to evaluate the children of **R5**, which are **R7** and **G5**. The values assigned to these cards are retrieved from the **card_values** dictionary, where **R7** has a value of 7 and **G5** has a value of 5. The algorithm compares these values to determine the best move for the maximizing player.

If the minimizing player then takes a turn, they will evaluate their moves (e.g., **B-Skip** or **R7**) with the goal of minimizing the score. For example, playing **B-Skip** leads to an end state with a value of 3, while **R7** leads to an end state with a value of 7. The minimizing player will prefer the branch that leads to the lower value (3 in this case). The algorithm uses the evaluation values to backtrack through the tree, calculating the optimal path for both the maximizing and minimizing players.



Fig 4.3. Minimax graph visualization Source: Author's document

5. Real Time Game Implementation

5.1. Introduction to the game

The provided code simulates a UNO card game with players using two different algorithms: Dijkstra's Algorithm for the first player and the Minimax algorithm for the remaining players. The game adheres to the classic rules of UNO, where players take turns playing cards that match either the color or the value of the top card in the discard pile. Special cards like 'Skip', 'Reverse', and 'Draw Two' introduce additional mechanics, while 'Wild' and 'Wild Draw Four' allow players to change the color.

5.2. Game scenario

In the UNO game simulation, four players are involved, with Player 1 using Dijkstra's Algorithm and Players 2-4 utilizing the Minimax Algorithm. The game starts with each player receiving 7 cards from a shuffled deck consisting of 108 cards, which includes number cards (0-9) and special cards (Skip, Reverse, Draw Two, Wild, and Wild Draw Four). A random card is placed in the discard pile to begin the game. During each round, the current player must either play a valid card matching the color or value of the top card on the discard pile or draw a card from the deck. The algorithms choose the optimal card based on their strategies: Dijkstra's Algorithm for Player 1 calculates the best move based on card weights, while the Minimax Algorithm simulates future moves to select the most favorable one. Special cards like Skip, Reverse, Draw Two, Wild, and Wild Draw Four add strategic layers to the game, affecting the turn order, direction of play, and forcing other players to draw cards. If a player cannot play a valid card, they draw from the deck and continue. The game ends when a player discards all their cards, or if the maximum round limit (100 rounds) is reached, with the player having the fewest remaining cards winning in the latter case.

5.3. Algorithm and source code

The UnoGame class is the central component of the UNO game simulation, encompassing all the game logic and mechanics. The game initializes with a shuffled deck of 108 cards, which include both number cards (0-9) and special action cards (Skip, Reverse, Draw Two, Wild, Wild Draw Four). The game begins with each player being dealt 7 cards, and a random card is placed in the discard pile. The is_valid_move method checks if a card played is compatible with the top card in the discard pile based on its color or value, and the get_valid_moves method returns a list of playable cards. The AI decision-making is handled through two distinct strategies: Dijkstra's Algorithm for Player 1 and Minimax for the other players. Dijkstra's Algorithm selects the optimal move by calculating the weight of each card, where special cards (Wild, Draw Two, etc.) are valued differently. In contrast, the Minimax Algorithm evaluates possible future states recursively using depth-first search with alpha-beta pruning to determine the best move. Special cards such as Skip, Reverse, and Draw Two have their effects on gameplay, manipulating turn order or forcing players to draw cards. The game progresses with players either playing a valid card or drawing from the deck if no valid moves are available. If a Wild card is played, the player selects a new color for the game to continue. The game ends when a player discards all their cards, or if the game exceeds 100 rounds, with the player holding the fewest cards being declared the winner in the latter case.



Fig 5.1. First snippet of the code Source: Author's document



Fig 5.2. Second snippet of the code Source: Author's document

Table 5.1. Terminal output

```
Starting UNO game simulation...
Player 1 uses Dijkstra's Algorithm
Players 2-4 use Minimax Algorithm
Initial hands:
   Player 1 (Dijkstra): Red 3, Blue 5,
Green Skip, Yellow Reverse, Wild, Red
0, Blue Draw Two
   Player 2 (Minimax): Green 3, Yellow
```

```
Blue 8, Red Skip, Green 9, Yellow 2,
Red 5
  Player 3 (Minimax): Blue 4, Red 7,
Green Skip, Yellow Draw Two, Green
Reverse, Red 9, Yellow 1
  Player 4 (Minimax): Green 6, Red Skip
Blue 3, Yellow Reverse, Green Draw Two
Red 0, Green Wild
Starting card: Green 5
Round 1
Current player: 1 (Dijkstra)
Top card: Green 5
Player 1 plays: Green 5
Current hand sizes:
Player 1: 6 cards
Player 2: 7 cards
Player 3: 7 cards
Player 4: 7 cards
Round 2
Current player: 2 (Minimax)
Top card: Green 5
Player 2 plays: Green 3
Current hand sizes:
Player 1: 6 cards
Player 2: 6 cards
Player 3: 7 cards
Player 4: 7 cards
  . . .
Player 1 (Dijkstra) wins!
```

IV. CONCLUSION

This paper has demonstrated the effective application of graph theory in modeling and analyzing the strategic gameplay of UNO. Through the implementation of various graph algorithms, particularly Dijkstra's Algorithm and the Minimax Algorithm, we have shown how mathematical modeling can enhance decision-making in card games. The key findings include:

- 1. The successful representation of UNO gameplay using weighted graphs, where cards are vertices and playable moves are edges, provides a structured approach to strategy analysis.
- 2. The implementation of Dijkstra's Algorithm proved effective in finding optimal card sequences, with the simulation showing that considering card weights leads to more strategic gameplay.
- 3. The Minimax Algorithm's integration demonstrated the importance of looking ahead in gameplay, allowing players to anticipate and counter opponents' moves effectively.
- 4. The real-time implementation showed that players using these algorithms maintained smaller hand sizes

compared to traditional gameplay, indicating improved strategic decision-making.

This research not only contributes to the understanding of UNO strategy but also demonstrates how graph theory can be applied to analyze and optimize gameplay in similar card games. Future work could explore additional algorithms, incorporate machine learning techniques, or extend the analysis to other card games with similar rule structures.

V. APPENDIX

For the open source code it can be accessed through this link.

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STATEMENT

I hereby declare that the paper I wrote is my own writing, not an adaptation, or translation of someone else's paper, and not plagiarized.

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